

# Deadlock

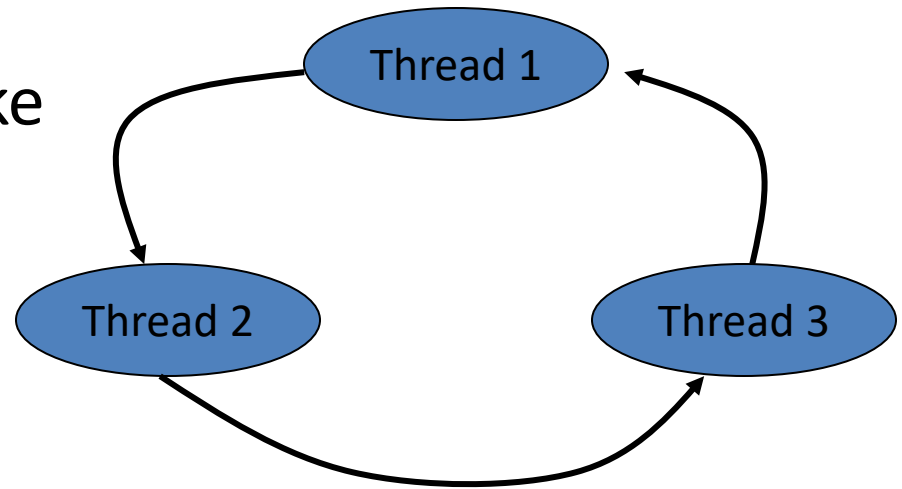
*การติดตาย*



<https://courses.cs.washington.edu/courses/cse451/06au/slides/os-deadlock.ppt>

# Deadlock

- Each thread is waiting on a resource held by another thread
  - So, there is no way to make progress

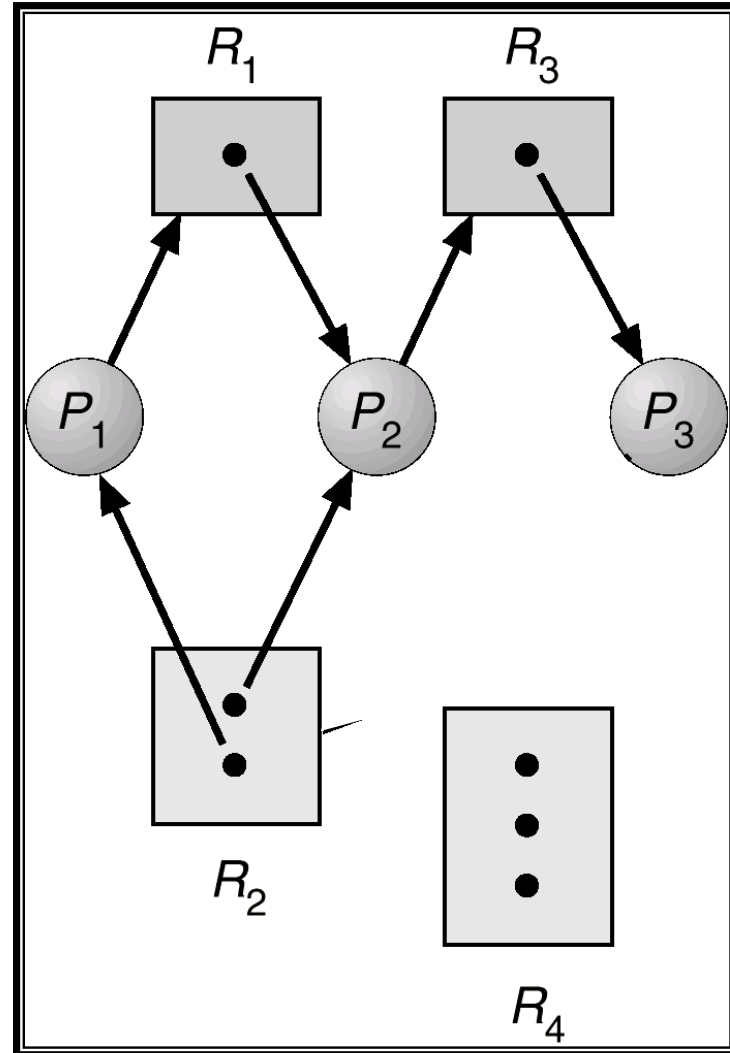


# Necessary Conditions for Deadlock

- **Mutual exclusion**
  - Resources cannot be shared
    - e.g., buffers, locks
- **Hold and wait (subset property)**
  - A thread must hold a *subset* of its resource needs
    - And, the thread is waiting for more resources
- **No preemption**
  - Resources cannot be taken away
- **Circular wait**
  - A needs a resource that B has
  - B has a resource that A has

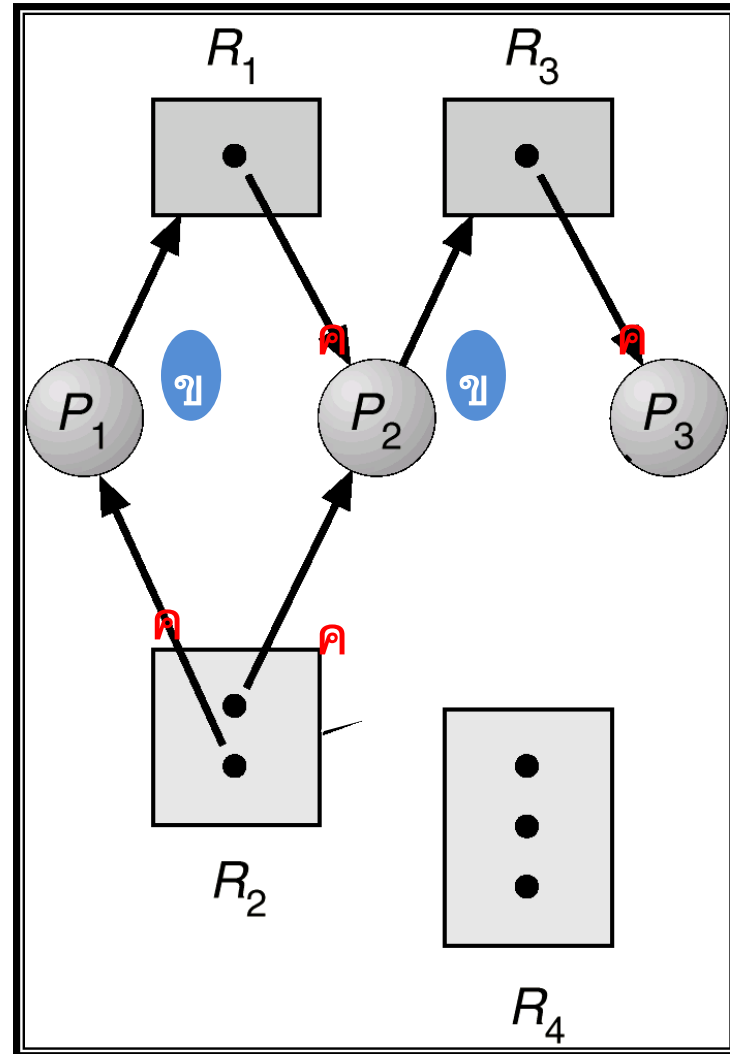
# Resource Allocation Graph **without** a Deadlock

$P \Rightarrow R$ : request edge  
 $R \Rightarrow P$ : assignment edge



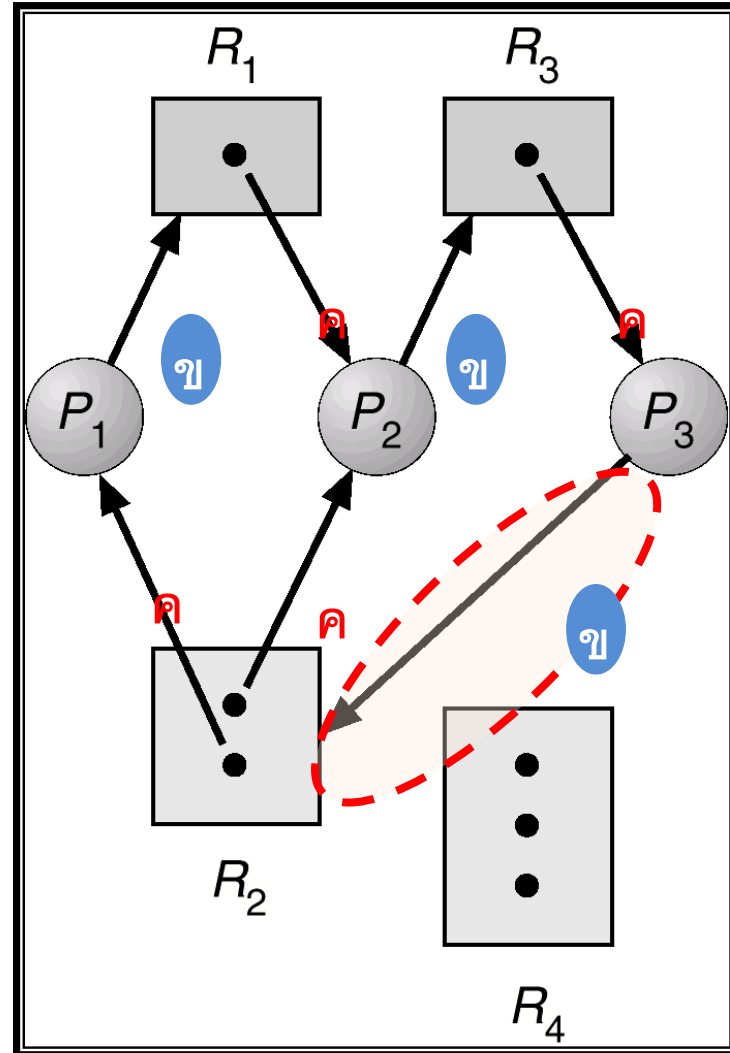
# Resource Allocation Graph **without** a Deadlock

$P \Rightarrow R$ : request edge  
 $R \Rightarrow P$ : assignment edge



# Resource Allocation Graph with a Deadlock

$P \Rightarrow R$ : request edge  
 $R \Rightarrow P$ : assignment edge



# Approaches to Deadlock

1. Avoid threads
2. Deadlock prevention
  - Break up one of the four necessary conditions
3. Deadlock avoidance
  - Stay live even in the presence of the four conditions
4. Detect and recover



# Deadlock Prevention

Can we eliminate:

- Mutual exclusion?
- Hold and wait (subset property)?
- No Preemption?
- Circular waiting?

# Deadlock Avoidance :

## Bankers Algorithm

- Basic idea: ensure that we always have an “escape route”
  - The resource graph is reducible
- This can be enforced with the bankers algorithm:
  - When a request is made
    - Pretend you granted it
    - Pretend all other legal requests were made
  - Can the graph be reduced?
    - If so, allocate the requested resource
    - If not, block the thread

# Deadlock Detection and Recovery

- Not commonly used
  - Detection is expensive
  - Recovery is tricky
- Possible exception: databases